Packing 4-cycles in Eulerian and bipartite Eulerian tournaments with an application to distances in interchange graphs

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Abstract

We prove that every Eulerian orientation of $K_{m,n}$ contains $\frac{1}{4+\sqrt{8}}mn(1-o(1))$ arc-disjoint directed 4-cycles, improving earlier lower bounds. Combined with a probabilistic argument, this result is used to prove that every regular tournament with n vertices contains $\frac{1}{8+\sqrt{32}}n^2(1-o(1))$ arc-disjoint directed 4-cycles. The result is also used to provide an upper bound for the distance between two antipodal vertices in interchange graphs.

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1 Introduction

All graphs and digraphs considered here are finite, and contain no parallel edges or anti-parallel arcs. For standard graph-theoretic terminology the reader is referred to [2]. An Eulerian orientation of an undirected graph G is an orientation of the edges of G such that each vertex has the same indegree and outdegree in the resulting oriented graph. It is well known that a graph G has an Eulerian orientation if and only if every vertex is of even degree. A tournament is an orientation of a complete graph and a bipartite tournament is an orientation of a complete bipartite graph. Properties of Eulerian tournaments (also called regular tournaments) and bipartite Eulerian tournaments have been extensively studied in the literature (see, e.g., [3, 8, 9]).

A conjecture of Brualdi and Shen [5] asserts that every bipartite Eulerian tournament has a decomposition into C_4 , the directed cycle on four vertices. In other words, any Eulerian orientation of $K_{m,n}$ contains mn/4 arc-disjoint C_4 's (clearly m and n must be even in order to have an Eulerian orientation). This conjecture is still far from being solved, and, in fact, there is no evidence that it should be true. Even an asymptotic version of this conjecture is not known. A recent result from

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[11] states that any Eulerian orientation of a bipartite graph with vertex class sizes m and n and more than 2mn/3 arcs has a C_4 . This immediately implies that any Eulerian orientation of $K_{m,n}$ contains at least mn/12 arc-disjoint C_4 's. Although very far from the Brualdi-Shen conjecture, this is currently the best known published lower bound for this problem. In this paper we considerably improve the lower bound.

Theorem 1.1 Every Eulerian orientation of $K_{m,n}$ contains at least $\frac{1}{4+\sqrt{8}}mn(1-o(1))$ arc-disjoint copies of C_4 .

The o(1) term denotes a function tending to zero as min $\{m, n\}$ tends to infinity. Notice that, trivially, any C_4 -packing of an orientation of $K_{m,n}$ contains at most mn/4 elements. Theorem 1.1 shows that more than 58 per cent of the arcs can be packed with C_4 's.

The motivation behind the Brualdi-Shen conjecture stems from the theory of interchange graphs. These graphs are defined as follows: Let $R = (r_1, \ldots, r_n)$ and $S = (s_1, \ldots, s_m)$ be non-negative integral vectors with $\sum r_i = \sum s_j$. Let $\mathcal{A}(R,S)$ denote the set of all $m \times n$ {0,1}-matrices with row sum vector R and column sum vector S, and assume that $\mathcal{A}(R,S) \neq \emptyset$. This set has been studied extensively (see [4] for a survey). An interchange is a transformation which replaces a 2×2 identity submatrix of a matrix A with the 2×2 permutation matrix $P_{i,j} = |i-j|$. Clearly an interchange (and hence any sequence of interchanges) does not alter the row and column sum vectors of a matrix and transforms a matrix in $\mathcal{A}(R,S)$ into another matrix in $\mathcal{A}(R,S)$. Ryser [10] proved that for any two matrices in $\mathcal{A}(R,S)$ there is a sequence of interchanges which transforms one to the other. The interchange graph G(R,S) of $\mathcal{A}(R,S)$, defined by Brualdi in 1980, is the graph with all matrices in $\mathcal{A}(R,S)$ as its vertices, where two matrices are adjacent provided that one can be obtained from the other by a single interchange. Brualdi conjectured that the diameter of G(R,S), denoted d(R,S), cannot exceed mn/4. This conjecture is still far from being resolved. The best known bounds for d(R, S) use the following reduction of Walkup. Given two matrices $A, B \in \mathcal{A}(R, S)$ consider the $m \times n$ matrix A - B. This matrix contains the elements $\{-1, 0, 1\}$ and the sum of each row and column is zero. Thus, there is a one-to-one correspondence between such matrices and Eulerian orientations of a bipartite graph with vertex class sizes m and n. Walkup [12] proved that the distance between A and B in G(R,S), denoted i(A,B) satisfies

$$i(A,B) = \frac{d(A,B)}{2} - q(A,B) \tag{1}$$

where d(A, B) is the number of nonzero entries in A - B and q(A, B) is the maximum number of arc-disjoint cycles in a cycle decomposition of the bipartite Eulerian digraph corresponding to A - B. The currently best known upper bound for d(R, S) obtained using Walkup's reduction is given in [11] yielding $d(R, S) \leq \frac{5}{12}mn$. We call two matrices $A, B \in \mathcal{A}(R, S)$ antipodal if B is the complement of A, or, in other words, if A + B is the all-one matrix. Notice that a necessary condition for A and B to be antipodal is that all coordinates of R are m/2 and all coordinates of

S are n/2. Trivially, the distance between two antipodal vertices in G(R,S) is at least mn/4, and hence such vertices show that Brualdi's conjecture, if true, is best possible. In particular antipodal vertices are conjectured to be farthest apart in G(R,S). If A and B are antipodal matrices then A-B corresponds to an Eulerian orientation of $K_{m,n}$. Observe that using (1) and Theorem 1.1 we obtain the following.

Corollary 1.2 Let A and B be two antipodal vertices of G(R,S). Then,

$$0.25mn \le i(A,B) \le \frac{1+\sqrt{2}}{4+\sqrt{8}}mn(1+o(1)) \approx 0.354mn(1+o(1)).$$

Theorem 1.1, combined with a probabilistic argument, enables us to prove an analogous theorem for regular tournaments.

Theorem 1.3 Every regular tournament with n vertices contains at least $\frac{1}{8+\sqrt{32}}n^2(1-o(1))$ arc-disjoint copies of C_4 .

Notice that, trivially, any C_4 -packing of an n-vertex tournament contains less than $n^2/8$ elements. Theorem 1.3 shows that more than 58 per cent of the arcs can be packed with C_4 's.

In the following section we prove Theorem 1.1. The final section contains the proof of Theorem 1.3.

2 Proof of Theorem 1.1

Kotzig [7] computed the minimum possible number of C_4 's in regular tournaments. We require an analogous lower bound for bipartite Eulerian tournaments. The following lemma provides this bound.

Lemma 2.1 Every bipartite Eulerian tournament with vertex class sizes m and n contains at least $m^3n^2/(32(m-1))$ copies of C_4 .

Proof: Let $G = (A \cup B, V)$ be a bipartite Eulerian tournament with vertex classes A and B where |A| = m and |B| = n. For each pair of distinct vertices $u, v \in A$, let $k_{u,v}$ denote the number of common out neighbors of u, v. Notice that $k_{u,v}$ is also the number of common incoming neighbors of u, v. Furthermore $n/2 - k_{u,v}$ is the number of out neighbors of u which are incoming neighbors of v and also the number of incoming neighbors of v which are out neighbors of v. The number of v copies containing both v and v is $(n/2 - k_{u,v})^2$. Thus, if v is the number of copies of v in v

then

$$x = \sum_{u,v \in A} \left(\frac{n}{2} - k_{u,v}\right)^{2}$$

$$= \frac{n^{2}}{4} \binom{m}{2} - n \sum_{u,v \in A} k_{u,v} + \sum_{u,v \in A} k_{u,v}^{2}$$

$$\geq \frac{n^{2}}{4} \binom{m}{2} - n \sum_{u,v \in A} k_{u,v} + \frac{1}{\binom{m}{2}} \left(\sum_{u,v \in A} k_{u,v}\right)^{2}$$

$$= \frac{n^{2}}{4} \binom{m}{2} - n^{2} \binom{m/2}{2} + \frac{n^{2}}{\binom{m}{2}} \binom{m/2}{2}^{2}$$

$$= \frac{n^{2}m^{3}}{32(m-1)}.$$

We note that Lemma 2.1 is asymptotically tight. This can easily be shown by considering random Eulerian orientations of $K_{m,n}$. We omit the details.

Our next lemma shows that if some arc of a bipartite Eulerian tournament is on too many copies of C_4 or on too few copies of C_4 then there must be many copies of C_4 altogether. Before we state the lemma we need a few definitions. Let G be a bipartite Eulerian tournament with vertex class sizes m and n. Let e = (x, y) be an arc of G. Let $N^-(x)$ denote the in-neighborhood of x and let $N^+(y)$ denote the out-neighborhood of x. Let x0 be the number of arcs from x1 to x2 to x3. Notice that x4 is on precisely x4 and notice also that there are x4 arcs from x5 to x6 to x7. Let x8 arcs from x7 to x8 to x9. Let x9 to x9 to x9 to x9 to x9 to x9. Let x9 to x9. Let x9 to x

Lemma 2.2 Let G be a bipartite Eulerian tournament with vertex class sizes m and n. G has at least $m^2n^2(1/16 - \alpha(G) + 4\alpha(G)^2)$ copies of C_4 .

Proof: Let A and B denote the vertex classes of G with |A| = m and |B| = n. Let e = (x, y) be an arc for which $\alpha(e) = \alpha(G)$. For simplicity let $\alpha = \alpha(e) = \alpha(G)$. Thus, either $\alpha = d(e)/mn$ or $\alpha = 1/4 - d(e)/mn$. Without loss of generality assume that $\alpha = 1/4 - d(e)/mn$ (the proof of the other case is analogous). Also without loss of generality assume that $x \in A$ and $y \in B$ (the proof of the other case is analogous). By our assumptions, $|N^-(x)| = |N^+(x)| = n/2$ and $|N^-(y)| = |N^+(y)| = m/2$. Also, e appears on precisely $d(e) = mn(1/4 - \alpha)$ copies of C_4 .

Let T be the set of arcs from $N^-(x)$ to $N^+(y)$. Notice that $|T| = \alpha mn$. Let S be the set of arcs from $N^+(x)$ to $N^-(y)$, let W be the set of arcs from $N^-(y)$ to $N^-(x)$ and let Z be the set of arcs from $N^+(y)$ to $N^+(x)$. Since G is Eulerian we must have $\alpha mn = |T| = |S| = |W| = |Z|$.

Consider an induced subgraph of G with four vertices, precisely one vertex from each of $N^+(x), N^-(x), N^+(y), N^-(y)$, and with no arc from $S \cup T \cup W \cup Z$. Such a subgraph must be a

 C_4 . Denote the number of such subgraphs by x. Denote the number of C_4 's of G having at least one arc in $S \cup T \cup W \cup Z$ by y. Denote the number of C_4 's of G with all their arcs in $S \cup T \cup W \cup Z$ by z. Clearly, $z \leq y$ and there are at least x + y copies of C_4 in G.

Let F be the set of induced subgraphs of G with four vertices, precisely one vertex from each of $N^+(x), N^-(x), N^+(y), N^-(y)$. Notice that $|F| = m^2 n^2 / 16$. Let $F' \subset F$ be the set of elements of F containing at least one arc from $S \cup T \cup W \cup Z$. Notice that all elements in $F \setminus F'$ are C_4 's. The sum of the arcs of $S \cup T \cup W \cup Z$ over all elements of F' is precisely $(4\alpha mn)(mn/4) = \alpha m^2 n^2$. However, some elements of F' contain more than one arc from $S \cup T \cup W \cup Z$, and are therefore counted more than once. In fact, each $e \in T$ and each $e' \in S$ appear together in precisely one element of F'. Similarly, each $e \in W$ and each $e' \in Z$ appear together in precisely one element of F'. Thus, there are $\alpha^2 m^2 n^2$ elements of F' containing a pair of arcs from $S \cup T$ and there are $\alpha^2 m^2 n^2$ elements of F' containing a pair of arcs from $F' \cap T$ and there are $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ and there are $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and there are $F' \cap T$ are some elements of $F' \cap T$ and the elements of $F' \cap T$ are some elements of $F' \cap T$ and the elements of $F' \cap T$ are some elements of $F' \cap T$ and the elements of $F' \cap T$ are some elements of $F' \cap T$ and $F' \cap T$ are some elements of

$$x = |F| - |F'| \ge \frac{1}{16}m^2n^2 - \alpha m^2n^2 + 2\alpha^2m^2n^2 + z.$$

By the definition of $\alpha = \alpha(G)$, each arc of G appears on at least αmn copies of C_4 . In particular the sum of the arcs from $S \cup T \cup W \cup Z$ appearing on copies of C_4 is at least $(4\alpha mn)(\alpha mn) = 4\alpha^2 m^2 n^2$. Precisely z of the C_4 's containing arcs from $S \cup T \cup W \cup Z$ are counted 4 times in this way. Notice that if some C_4 has three arcs from $S \cup T \cup W \cup Z$ then it must have four arcs in $S \cup T \cup W \cup Z$. Thus, $y \geq z + (4\alpha^2 m^2 n^2 - 4z)/2 = 2\alpha^2 m^2 n^2 - z$. Consequently,

$$x + y \ge m^2 n^2 (\frac{1}{16} - \alpha + 4\alpha^2).$$

Recall that a k-uniform hypergraph is a family of k-subsets of some set V of vertices. The k-subsets are the edges of the hypergraph. If x, y are two vertices of a hypergraph then let d(x) denote the number of edges containing x and let d(x, y) denote the number of edges containing both x and y (their co-degree). A proper r-coloring of the edges of a hypergraph is a partition of the set of edges to r subsets such that any two edges in the same subset are disjoint. Kahn [6] has proved a very powerful result giving an upper bound for the minimum number of colors in a proper edge-coloring of a uniform hypergraph (his result is, in fact, more general).

Lemma 2.3 (Kahn [6]) For every $k \geq 2$ and every $\epsilon > 0$ there exists a $\delta > 0$ such that the following statement is true:

If H is a k-uniform hypergraph on a set V of vertices, and if $d(v) \leq D$ for all $v \in V$ and $d(u,v) < \delta D$ for all distinct $u,v \in V$, then there is a proper coloring of the edges of H with at most $(1+\epsilon)D$ colors.

Completing the proof of Theorem 1.1. Let $\epsilon > 0$. We need to prove that if m is sufficiently large as a function of ϵ and if $n \geq m$ and G is a bipartite Eulerian tournament with vertex class sizes m and n, then G has at least $\frac{1}{4+\sqrt{8}}mn(1-\epsilon)$ arc-disjoint copies of C_4 .

We construct a 4-uniform hypergraph H as follows. The vertices of H are the mn arcs of G and the edges of H are the copies of C_4 in G. Let e = (x,y) be an arc of G such that d(e) is maximal. Clearly, $d(e) \leq mn/4$. By Lemma 2.1 we also have $d(e) \geq 4(m^2n^2/32)/(mn) \geq mn/8$. In particular $\alpha(e) = 1/4 - d(e)/mn$ and $0 \leq \alpha(G) \leq \alpha(e) \leq 1/8$. Notice that any two arcs of G appear together on at most $\max\{m/2, n/2\} = n/2$ copies of C_4 . Thus, with D = d(e) and with $\delta = \delta(4, \epsilon)$ chosen as in Lemma 2.3 we have that for m sufficiently large, the hypergraph H satisfies the conditions of Lemma 2.3. In particular, there is a coloring of the C_4 copies of G with at most $(1 + \epsilon)d(e)$ colors such that each color class corresponds to a set of arc-disjoint copies of C_4 . By Lemma 2.1 and Lemma 2.2, the total number of edges of H is at least

$$\max\{\frac{1}{32}m^2n^2, m^2n^2(\frac{1}{16}-\alpha(G)+4\alpha(G)^2)\}.$$

Thus, there must be a set of arc-disjoint copies of C_4 with at least

$$\max\{\frac{m^2n^2/32}{(1+\epsilon)d(e)}, \frac{m^2n^2(1/16-\alpha(G)+4\alpha(G)^2)}{(1+\epsilon)d(e)}\}$$

elements. Since $\alpha(G) \leq \alpha \leq 1/8$ it suffices to prove that

$$\frac{mn}{1+\epsilon} \max\{\frac{1/32}{1/4 - \alpha(e)}, \frac{1/16 - \alpha(e) + 4\alpha(e)^2}{1/4 - \alpha(e)}\} \ge \frac{1}{4+\sqrt{8}} mn(1-\epsilon).$$

Indeed, this follows from the fact that for $0 \le z \le 1/8$ the value of $\max\{(1/32)/(1/4-z)$, $(1/16-z+4z^2)/(1/4-z)\}$ is minimized when $z=1/8-\sqrt{1/128}$.

3 Proof of Theorem 1.3

The proofs of Lemma 2.1 and Lemma 2.2, and hence of Theorem 1.1, can easily be modified to hold for almost Eulerian bipartite tournaments. More precisely, we say that a bipartite orientation of $K_{m,n}$ is δ -Eulerian if the outdegree and indegree of each vertex from the m-vertex class is at least $(1 - \delta)n/2$ and if the outdegree and indegree of each vertex from the n-vertex class is at least $(1 - \delta)m/2$. Given $\epsilon > 0$, it is straightforward to check that Lemma 2.1 holds for δ -Eulerian bipartite tournaments if one replaces the $m^2n^2/32$ bound with $m^2n^2(1-\epsilon)/32$ provided that $\delta = \delta(\epsilon)$ is sufficiently small and min $\{m,n\}$ is sufficiently large as a function of ϵ . It is straightforward to check that a similar relaxation also holds for Lemma 2.2. Thus, we obtain the following relaxation for Theorem 1.1.

Lemma 3.1 Let $\epsilon > 0$. There exists $\delta = \delta(\epsilon) > 0$ and $M = M(\epsilon)$ such that if $M \le m \le n$, any δ -Eulerian orientation of $K_{m,n}$ contains at least $\frac{1}{4+\sqrt{8}}mn(1-\epsilon)$ arc-disjoint copies of C_4 .

Completing the proof of Theorem 1.3. Let $\epsilon > 0$. We need to show that there exists $N = N(\epsilon)$ such that for all n > N, any regular tournament with n vertices has at least $\frac{1}{8+\sqrt{32}}n^2(1-\epsilon)$ arc-disjoint directed 4-cycles.

We choose $\delta = \delta(\epsilon/2)$ to be as in Lemma 3.1. Choose $N \ge \max\{2M^2, 64/\epsilon^2\}$, where $M = M(\epsilon/2)$ is the constant from Lemma 3.1, to be the smallest integer satisfying $4N^2 \exp(-\delta^2 N/128) < 1$. Let n > N. The asymptotic nature of our proof allows us to assume, for simplicity, that n is a square of an integer. Therefore, let $m = \sqrt{n}$. Let T be a regular tournament with n vertices. Each $v \in T$ chooses uniformly and independently an integer from $\{1, \ldots, m\}$. Let V_i denote the set of vertices that chose i, and for $1 \le i < j \le m$ let $G_{i,j}$ be the complete bipartite tournament with vertex classes V_i and V_j .

We will prove that with positive probability, all the $\binom{m}{2}$ bipartite tournaments $G_{i,j}$ are δ -Eulerian, and that all V_i have size at least M and at most 2m. This suffices to complete the theorem since by Lemma 3.1 we have that the total number of arc-disjoint C_4 in T is at least

$$\sum_{1 \le i < j \le m} \frac{1}{4 + \sqrt{8}} |V_i| |V_j| (1 - \frac{\epsilon}{2})$$

$$= \frac{1}{8 + \sqrt{32}} (1 - \frac{\epsilon}{2}) (n^2 - \sum_{i=1}^m |V_i|^2) \ge \frac{1}{8 + \sqrt{32}} (1 - \frac{\epsilon}{2}) (n^2 - 4m^3) \ge$$

$$\ge \frac{1}{8 + \sqrt{32}} (1 - \frac{\epsilon}{2}) n^2 (1 - \frac{\epsilon}{2}) \ge \frac{1}{8 + \sqrt{32}} n^2 (1 - \epsilon).$$

Fix a vertex v. Let $d_i^+(v)$ (resp. $d_i^-(v)$) denote the number of out (resp. incoming) neighbors of v in V_i . The expectation of $d_i^+(v)$ (resp. $d_i^-(v)$) is precisely $m^{-1}(n-1)/2 = (n-1)/(2m)$. A large deviation inequality of Chernoff (cf. [1] Appendix A) bounds the probability that a random variable which is the sum of independent indicator random variables deviates from its expectation by more than a specified amount. More precisely, for the variable $d_i^+(v)$ and the amount $\frac{\delta}{4} \frac{n-1}{2m}$ we get

$$\Pr\left[\left|d_i^+(v) - \frac{n-1}{2m}\right| > \frac{\delta}{4} \frac{n-1}{2m}\right] < 2\exp\left(-\frac{\delta^2(n-1)^2/64m^2}{2(n-1)/2m} + \frac{\delta^3(n-1)^3/512m^3}{2(n-1)^2/4m^2}\right) = 2\exp\left(-\frac{\delta^2(n-1)}{64m} + \frac{\delta^3(n-1)}{128m}\right) < \frac{1}{2mn}.$$

Precisely the same bound holds for $d_i^-(v)$. Since there are n choices for vertices and m choices for i, we get that with probability greater than $1 - 2 \cdot n \cdot m \cdot \frac{1}{2mn} = 0$, for all vertices v and all indices i,

$$\left|d_i^+(v) - \frac{n-1}{2m}\right| \le \frac{\delta}{4} \frac{n-1}{2m} \qquad \left|d_i^-(v) - \frac{n-1}{2m}\right| \le \frac{\delta}{4} \frac{n-1}{2m}.$$

We therefore fix a partition V_1, \ldots, V_m for which all of these conditions are met. In particular, this means that for all $i = 1, \ldots, m$,

$$2m > (1 + \frac{\delta}{4})\frac{n-1}{m} \ge |V_i| \ge (1 - \frac{\delta}{4})\frac{n-1}{m} > M,$$

and that for all vertices v and all i = 1, ..., m

$$\min\{d_i^-(v) , d_i^+(v)\} \ge (1 - \frac{\delta}{4}) \frac{n-1}{2m} \ge (1 - \delta)(1 + \frac{\delta}{4}) \frac{n-1}{2m} \ge \frac{|V_i|}{2}(1 - \delta).$$

Thus, all the $\binom{m}{2}$ bipartite tournaments $G_{i,j}$ are δ -Eulerian, and all V_i have size at least M and at most 2m.

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