

Sum-distinguishing number of sparse hypergraphs

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Abstract

A vertex labeling of a hypergraph is *sum distinguishing* if it uses positive integers and the sums of labels taken over the distinct hyperedges are distinct. Let $s(H)$ be the smallest integer N such that there is a sum-distinguishing labeling of H with each label at most N . The largest value of $s(H)$ over all hypergraphs on n vertices and m hyperedges is denoted $s(n, m)$. We prove that $s(n, m)$ is almost-quadratic in m as long as m is not too large. More precisely, the following holds: If $n \leq m \leq n^{O(1)}$ then

$$s(n, m) = \frac{m^2}{w(m)},$$

where $w(m)$ is a function that goes to infinity and is smaller than any polynomial in m .

The parameter $s(n, m)$ has close connections to several other graph and hypergraph functions, such as the irregularity strength of hypergraphs. Our result has several applications, notably:

- We answer a question of Gyárfás et al. whether there are n -vertex hypergraphs with irregularity strength greater than $2n$. In fact we show that there are n -vertex hypergraphs with irregularity strength at least $n^{2-o(1)}$.
- In addition, our results imply that $s^*(n) = n^2/w(n)$ where $s^*(n)$ is the distinguishing closed-neighborhood number, i.e., the smallest integer N such that any n -vertex graph allows for a vertex labeling with positive integers at most N so that the sums of labels on distinct closed neighborhoods of vertices are distinct.

1 Introduction

For a hypergraph $H = (V, E)$, we say that a labeling $f : V \rightarrow \mathbb{N}$ is *sum-distinguishing* or simply *distinguishing* if $s(e) \neq s(e')$ for any two distinct hyperedges $e, e' \in E$, where $s(e) = \sum_{v \in e} f(v)$. Let $s(H)$ be the smallest integer N such that there is a distinguishing labeling of H with each label at most N . Note that $s(H)$ is well-defined by assigning vertex labels equal to distinct powers of 2. Distinguishing labelings can be viewed as number-theoretic constructions extending Sidon sets to non-complete, non-uniform hypergraphs. Using common notation, a $B_h[1]$ -Sidon set is a set X of integers such that for any integer q , there is at most one subset X' of X , $|X'| = h$, so that the sum

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of elements from X' is q . So, a $B_h[1]$ -Sidon set corresponds to an injective distinguishing labeling of a complete h -uniform hypergraph. On the other hand, distinguishing labelings of hypergraphs are closely connected to several “distinguishing” type parameters of graphs and hypergraphs that we discuss in more detail later. Let

$$s(n, m) = \max\{s(H) : |V(H)| = n, |E(H)| = m\}.$$

Namely, $s(n, m)$ is the largest value of $s(H)$ over all hypergraphs on n vertices and m hyperedges. Observe first that for the largest possible value of m , namely $m = 2^n - 1$ (corresponding to the full hypergraph consisting of all possible hyperedges), it trivially holds that $s(n, 2^n - 1) \leq 2^{n-1} \leq m$. So, in particular, we have that $s(n, m)$ is at most linear in the number of edges whenever $m = \Theta(2^n)$. On the other hand, for general m , a standard probabilistic argument shows that $s(n, m) = O(m^2)$. So, it seems of interest to study the dependence of $s(n, m)$ on m whenever the hypergraph is relatively sparse. Our main result does just that. We prove, perhaps surprisingly, that for hypergraphs with polynomially many edges, $s(n, m)$ is neither linear nor quadratic. In fact, we prove that in this regime, $s(n, m)$ is almost-quadratic in m .

Theorem 1. *If $n \leq m \leq n^{O(1)}$ then*

$$s(n, m) = \frac{m^2}{w(m)},$$

where $w(m)$ is a function that goes to infinity and is smaller than any polynomial in m . More formally, for any $C > 0$, $\epsilon > 0$, there is n_0 such that for any $n > n_0$, and any m satisfying $n \leq m \leq n^{1/\epsilon}$, we have that $m^{2-\epsilon} \leq s(n, m) \leq m^2/C$.

The upper bound in the proof of Theorem 1 relies on several probabilistic arguments, some of which are rather delicate. For the lower bound, we extend an approach of Bollobás and Pikhurko [8] used for 2-uniform hypergraphs (i.e. graphs) and their distinguishing labelings.

Our main result has several applications that we next describe. Our first application is to the problem of distinguishing the vertices of a graph by sums of labels on closed neighborhoods. For a graph $G = (V, E)$, and a vertex $v \in V(G)$, the open neighborhood of v is $N(v) = \{u \in V(G) : uv \in E(G)\}$; the closed neighborhood of v is $N[v] = \{v\} \cup N(v)$. For a vertex labeling f of G and $v \in V(G)$, let $s^*(v) = s_f^*(v) = \sum_{x \in N[v]} f(x)$. The labeling f is called *vertex sum-distinguishing* if it uses positive integers and $s^*(v) \neq s^*(u)$ for any $u, v \in V(G)$ such that $N[u] \neq N[v]$. Let $s^*(G)$ be the smallest integer k such that there is a vertex sum-distinguishing labeling of G with a largest label k and let $s^*(n)$ be the maximum of $s^*(G)$ taken over all graphs with n vertices.

Let $s(n) = s(n, n)$. We observe that the parameters $s(n)$ and $s^*(n)$ are closely connected. Indeed, for a graph $G = (V, E)$ consider a hypergraph $H = H_G$ on a vertex set V with hyperedges corresponding to the closed neighbourhoods of vertices of G . We see that $s^*(G) = s(H)$. Note that the number of hyperedges in H_G is at most n . The following result is an immediate consequence of Theorem 1 and Lemma 1, in which we prove that $s(n/2) \leq s^*(n) \leq s(n)$. Thus, we obtain:

Corollary 1. *We have that*

$$s^*(n) = \frac{n^2}{w(n)},$$

where $w(n)$ is a function that goes to infinity and is smaller than any polynomial in n . More formally, for any $C > 0$, $\epsilon > 0$, there is n_0 such that for any $n > n_0$, $n^{2-\epsilon} \leq s^*(n) \leq n^2/C$.

The proof of Theorem 1 (and hence Corollary 1) yields an efficient randomized algorithm for finding a corresponding labeling. For graphs with given maximum and minimum degrees, we provide a more specific result which also yields an efficient deterministic algorithm.

Theorem 2. *Let G be a nonempty n -vertex graph with maximum degree Δ , minimum degree δ , and the largest number of vertices with pairwise distinct closed neighborhoods equal to n' . Let $d(v)$ denote the degree of v . Then*

$$\frac{n' + \delta}{\Delta + 1} \leq s^*(G) \leq \max\{(n - d(v) - 1)(d(v) + 1) + 2 : v \in V(G)\} \leq (\Delta + 1)n.$$

Observe that the upper bound of Theorem 2 is weaker than the upper bound in Corollary 1 whenever $\Delta = \Theta(n)$. In fact, it only gives $s^*(n) \leq n^2/4 + 2$. As a final result concerning s^* , we consider the case where the graph is a tree. For a tree T and its vertex u , let $L(u)$ be the set of leaves adjacent to u . Let $L(T) = \max\{|L(u)| : u \in V(T)\}$.

Theorem 3. *Let T be a tree with $n \geq 3$ vertices. Then $s^*(T) \leq 2n - 2 - L(T)$, moreover this bound is tight for stars.*

The parameter $s(n)$ is closely related to the notion of *irregularity strength* of hypergraphs. There is extensive literature on irregularity strengths of graphs, a notion first introduced (for graphs) by Chartrand et al. [9], see also for example Nierhoff [13] and Blokhuis and Szőnyi [7], Balister et al. [4], as well as the survey by Gallian [10]. To define irregularity strength, consider an edge-labeling f of a hypergraph H with positive integers and for each vertex x compute $s(x)$, the sum of labels over all hyperedges containing x . The labeling is *irregular*, if the sums s are distinct for all vertices. The smallest value of a largest label used in an irregular labeling of H is denoted $irr(H)$ and $irr(n)$ is the largest value of $irr(H)$ over all n -vertex hypergraphs. Note that $irr(H)$ corresponds to $s(H^*)$, where H^* is the dual hypergraph of H . Recall that for a hypergraph $H = (V, E)$, the dual hypergraph H^* has vertex set E and edge set $\{e \in E : e \ni x\} : x \in V$. Gyárfás et al. [11] provided upper bounds on $irr(H)$ and stated that “it is not known whether $irr(n) \geq 2n$.” A consequence of our result gives a better lower bound $irr(n) \geq n^{2-\epsilon}$ for any positive ϵ and sufficiently large n , and in particular, answers their question.

Theorem 4. *For any $\epsilon > 0$, there is n_0 such that for any $n > n_0$, $irr(n) \geq n^{2-\epsilon}$.*

We mention a few other closely related problems that have been studied. There is yet another parameter, similar to $s(H)$, introduced by Bhattacharya et al. [6] and called a *discriminator* where the goal is to assign non-negative integer labels to vertices of a hypergraph such that the sums on the hyperedges are distinct, and positive. While our original motivation was to distinguish

the vertices of a graph via sums on closed neighborhoods, there is a similar problem restricted to pairs of vertices that are adjacent, i.e., so-called adjacent vertex sum-distinguishing number, that was studied for closed neighborhoods by Axenovich et al. [3] and for open neighborhoods by Bartnicki et al. [5], who use also an unpublished observation by Norin. These above-mentioned adjacency-dependent parameters can however be upper-bounded by a function of the maximum degree, independent of the number of vertices of the graph. Finally, we mention that distinguishing labelings of *graphs* were also studied by Ahmad et al. [1].

The rest of the paper is structured as follows. In the next section we prove several lemmas that are required for our theorems. In particular, Lemma 1 comparing $s^*(n)$ and $s(n)$, Lemma 2, which is the main ingredient in the lower bound on $s(n, m)$ as it implies the existence of a certain (randomly constructed) hypergraph H with large $s(H)$, and Lemma 3 about the distribution of the sum of discrete random variables, that we use for the upper bound on $s(n, m)$. In Section 3 we prove Theorem 1, our main result. Section 4 contains the proofs of Theorems 2, 3, and 4. The final section consists of concluding remarks and open problems.

2 Lemmas

This section consists of several lemmas facilitating the proof of our main theorems. For a positive integer x , we use the notation $[x] = \{1, \dots, x\}$. Our first lemma relates $s^*(n)$, $s(n)$, and $s^*(2n)$.

Lemma 1. *For any $n \geq 2$, we have $s^*(n) \leq s(n) \leq s^*(2n)$.*

Proof. Let G be an n -vertex graph with $s^*(G) = s^*(n)$. As mentioned in the introduction, consider a hypergraph H on the vertex set $V = V(G)$ with hyperedges corresponding to the closed neighborhoods of vertices in G . Since a labeling f of V is vertex sum-distinguishing in G if and only if it is distinguishing in H , we have that $s^*(n) = s^*(G) = s(H) \leq s(n)$ ¹. On the other hand, consider a hypergraph H on a vertex set $B = \{b_1, \dots, b_n\}$ and with n hyperedges e_1, \dots, e_n , such that $s(H) = s(n)$. Let G be a graph on vertex set $A \cup B$, where $A = \{a_1, \dots, a_n\}$, $A \cap B = \emptyset$, where A induces a clique with n vertices, B induces an independent set, and $a_i b_j \in E(G)$ if and only if $b_j \in e_i$. Then we see that if a labeling f is vertex sum-distinguishing in G then, restricted to B , it is distinguishing in H . Consider such an optimal f , i.e. with a largest label $s^*(G)$. Since G has $2n$ vertices, $s^*(G) \leq s^*(2n)$. Thus, $s(n) = s(H) \leq s^*(G) \leq s^*(2n)$. \square

Lemma 2. *For any fixed $r \geq 2$, there is a constant $c = c(r)$ such that for every positive integer N there exists an r -uniform hypergraph H on N vertices such that*

$$\begin{aligned} |E(H)| &= \Theta(N^{(r+1)/2} \sqrt{\log N}), \text{ and} \\ s(H) &\geq cN^r. \end{aligned}$$

Proof. We are going to extend a result of Bollobás and Pikhurko [8] on distinguishing labelings of graphs to r -uniform hypergraphs. Also note that the inequalities in the lemma's statement allow

¹Observe that H might have less than n edges since not all closed neighborhoods of G are necessarily distinct, but since adding edges to a hypergraph cannot decrease s , we indeed have $s(H) \leq s(n)$

us to assume, whenever necessary, that N is sufficiently large as a function of r .

Proof idea: We provide a lower bound on $s(H)$ for a random r -uniform hypergraph $H \sim G_r(N, p)$, i.e., a hypergraph on a vertex set $[N]$, such that hyperedges are chosen independently with probability p . In order to show that $s(H) > s$ for a chosen s , we shall consider a fixed labeling f of $[N]$ and denote by p' the probability that f is distinguishing for H . Now, if it holds that $p' = o(s^{-N})$ then we have $\Pr[s(H) \leq s] \leq s^N p' = o(1)$. So, in this case we see that almost surely $s(H) > s$.

Let $q = \sqrt{13r \cdot r!}$, $p = q\sqrt{\ln N}/\sqrt{N^{r-1}}$, $s = \lfloor N^r/(2r \cdot r!) \rfloor$ and $H \sim G_r(N, p)$. Consider a labeling $f : [N] \rightarrow [s]$. For any $e \in \binom{[N]}{r}$, let $s(e) = \sum_{i \in e} f(i)$. We estimate p' , the probability that f is distinguishing for H .

Let H_k be the r -uniform hypergraph on a vertex set $[N]$, with $E(H_k) = \{e \in \binom{[N]}{r} : s(e) = k\}$ and denote $h_k = |E(H_k)|$. Note that for any r -subset of the vertices $e \in \binom{[N]}{r}$, $r \leq s(e) \leq sr$ and that the H_k 's form an edge-decomposition of the complete r -uniform hypergraph on the vertex set $[N]$. Note that f is distinguishing for H if and only if H has at most one edge in each of the H_k 's. We need to consider only those H_k 's that have at least two edges so let $K = \{k : h_k \geq 2\}$. We have

$$\begin{aligned} p' &= \Pr[f \text{ is distinguishing for } H] \\ &= \prod_{k \in K} \Pr[|E(H) \cap E(H_k)| \leq 1] \\ &= \prod_{k \in K} \left((1-p)^{h_k} + h_k p (1-p)^{h_k-1} \right). \end{aligned}$$

We need the following statement that is a routine calculation. If $t_1 \leq t_2 - 2$ then

$$\begin{aligned} & \left((1-p)^{t_1} + t_1 p (1-p)^{t_1-1} \right) \left((1-p)^{t_2} + t_2 p (1-p)^{t_2-1} \right) \\ & \leq \left((1-p)^{t_1+1} + (t_1+1)p(1-p)^{(t_1+1)-1} \right) \left((1-p)^{t_2-1} + (t_2-1)p(1-p)^{(t_2-1)-1} \right). \quad (1) \end{aligned}$$

Using (1) we can upper-bound the expression for p' by the one in which each h_k takes an integer value x or $x+1$, for some x . Let there be b of x 's and $|K| - b$ of $(x+1)$'s, so $bx + (|K| - b)(x+1) = \sum_{k \in K} h_k = h$. Assume that $xb \geq h/2$ (the case where $(|K| - b)(x+1) \geq h/2$ is analogous). We have:

$$\begin{aligned} p' &= \prod_{k \in K} (1-p)^{h_k} + h_k p (1-p)^{h_k-1} \\ &\stackrel{(1)}{\leq} \left((1-p)^x + xp(1-p)^{x-1} \right)^b \left((1-p)^{x+1} + (x+1)p(1-p)^{(x+1)-1} \right)^{|K|-b} \\ &\leq \left((1-p)^x + xp(1-p)^{x-1} \right)^b \\ &\leq \left((1-p)^x + xp(1-p)^{x-1} \right)^{\frac{h}{2}/x}. \end{aligned}$$

It is also a routine calculation, that for any any p , $0 < p < 1$

$$\max_{t \geq 2} ((1-p)^t + tp(1-p)^{t-1})^{1/t} \stackrel{t=2}{=} ((1-p)^2 + 2p(1-p))^{1/2}. \quad (2)$$

Coming back to bounding p' , we have

$$\begin{aligned} p' &\leq ((1-p)^x + xp(1-p)^{x-1})^{\frac{h}{2}/x} \\ &\stackrel{(2)}{\leq} ((1-p)^2 + 2p(1-p))^{h/4} \\ &= (1-p^2)^{h/4} \\ &\leq e^{-p^2 h/4}. \end{aligned}$$

Observe also that the total number of hypergraphs H_k is at most

$$sr \leq \frac{N^r}{2r!} \leq \frac{1}{2} \binom{N}{r} (1 + o(1)).$$

Thus, at least about a half of the possible r -sets of vertices belong to H_k 's that have at least two edges, i.e., to H_k 's, $k \in K$. In other words, for sufficiently large N ,

$$h = \sum_{k \in K} h_k \geq \binom{N}{r} - \frac{1}{2} \binom{N}{r} (1 + o(1)) \geq \frac{N^r}{3r!}.$$

Recall that $p = q\sqrt{\ln N}/\sqrt{N^{r-1}}$ and $q = \sqrt{13r \cdot r!}$. Then

$$\begin{aligned} p' &\leq e^{-p^2 h/4} \\ &\leq e^{-q^2 (\ln N / N^{r-1}) N^r / (12r!)} \\ &= e^{-(q^2 / 12r!) N \ln N} \\ &= e^{-(13r/12) N \ln N} \\ &\leq N^{-rN}. \end{aligned}$$

Also recall that $s = \lfloor N^r / (2r \cdot r!) \rfloor$ so,

$$\begin{aligned} \Pr[s(H) \leq s] &\leq s^N p' \\ &\leq s^N N^{-rN} \\ &= o(1). \end{aligned}$$

This implies that with high probability $s(H) > s = cN^r$, for a constant c depending on r . Moreover, with high probability $|E(H)| = \Theta(p \binom{N}{r}) = \Theta(pN^r) = \Theta(N^{(r+1)/2} \sqrt{\log N})$. \square

Our final lemma of this section upper-bounds the probability that a sum of i.i.d. uniform discrete random variables attains a particular value. We will use it as an ingredient in the upper

bound proof of Theorem 1.

Lemma 3. *For any constant $C > 0$, there exists ℓ_0 such that for any $\ell > \ell_0$ the following holds. There is an integer $N_0 = N_0(\ell)$ such that for any $N > N_0$, if $X_1, \dots, X_{2\ell}$ are i.i.d. uniform discrete random variables over $[N]$, then for any integer t , $\Pr[X_1 + \dots + X_{2\ell} = t] \leq 5/(e^{4C} N)$.*

Proof. We first establish the following claim, asserting the concavity of a sum of i.i.d. uniform discrete random variables.

Claim 1. *Let X_1, \dots, X_ℓ be i.i.d. uniform discrete random variables over $[N]$ and let $X = X_1 + \dots + X_\ell$. Then for every real $d \geq 0$ it holds that*

$$\Pr[X = (N + 1)\ell/2 - d] = \Pr[X = (N + 1)\ell/2 + d] \geq \Pr[X = (N + 1)\ell/2 + d + 1].$$

Proof. It will be slightly more convenient to define $W_i = X_i - 1$, $i = 1, \dots, \ell$, $W = W_1 + \dots + W_\ell$, $q = N - 1$ and prove the equivalent statement

$$\Pr[W = q\ell/2 - d] = \Pr[W = q\ell/2 + d] \geq \Pr[W = q\ell/2 + d + 1].$$

Observe first that $\Pr[W = q\ell/2 - d] = \Pr[W = q\ell/2 + d]$ as W is symmetric around its mean $q\ell/2$. Observe next that $W \in \{0, \dots, q\ell\}$ so the inequality is only interesting if $d + 1 \leq q\ell/2$ and $q\ell/2 + d$ is an integer. We prove the claim by induction on ℓ . The case $\ell = 1$ trivially holds as W_1 is uniform. Assuming the claim holds for $\ell - 1$, we prove it for ℓ . Let $W^* = W_1 + \dots + W_{\ell-1}$ so $W = W^* + W_\ell$. We prove that $\Pr[W = q\ell/2 + d] \geq \Pr[W = q\ell/2 + d + 1]$. Observe that

$$\Pr[W = q\ell/2 + d] = \sum_{k=0}^q \Pr[W^* = q\ell/2 + d - k] \cdot \Pr[W_\ell = k] = \frac{1}{q+1} \sum_{k=0}^q \Pr[W^* = q\ell/2 + d - k].$$

Similarly,

$$\Pr[W = q\ell/2 + d + 1] = \frac{1}{q+1} \sum_{k=0}^q \Pr[W^* = q\ell/2 + d + 1 - k].$$

So,

$$\begin{aligned} & (q+1) (\Pr[W = q\ell/2 + d] - \Pr[W = q\ell/2 + d + 1]) \\ &= \Pr[W^* = q\ell/2 + d - q] - \Pr[W^* = q\ell/2 + d + 1] \\ &= \Pr[W^* = q(\ell-1)/2 + d - q/2] - \Pr[W^* = q(\ell-1)/2 + d + 1 + q/2] \\ &\geq 0 \end{aligned}$$

where the last inequality follows from the induction hypothesis and from the fact that $|d - q/2| < |d + 1 + q/2|$. \square

To prove the lemma, let

$$f_1 = X_1 + \dots + X_\ell \quad \text{and} \quad f_2 = X_{\ell+1} + \dots + X_{2\ell}.$$

We estimate the probability $\Pr[f_1 + f_2 = t]$. We will first need to prove two additional claims. The first is an anti-concentration result for f_1, f_2 and the second is a concentration result for them. Throughout the remainder of the proof we assume that ℓ is sufficiently large as a function of C and that N is sufficiently large as a function of ℓ .

Claim 2. *Let $j \in \{1, 2\}$. For any $C > 0$, there exists $\gamma = \gamma(C) > 0$ such that for every real number x it holds that*

$$\Pr[x - \gamma\sqrt{\ell}N \leq f_j \leq x + \gamma\sqrt{\ell}N] \leq e^{-4C} . \quad (3)$$

Proof. Recall that each X_i is uniform discrete over $[N]$. For the sake of our analysis it would be convenient to obtain X_i as follows. Let $U_i \sim U[0, 1]$ (i.e. U_i is uniform continuous in $[0, 1]$). Define $X_i = \lceil NU_i \rceil$. Since $\Pr[U_i = 0] = 0$, we have that X_i is discrete uniform over $[N]$ as the probability that $X_i = t$ is $1/N$ for each $t \in [N]$. Denote $g_1 = U_1 + \dots + U_\ell$ and $g_2 = U_{\ell+1} + \dots + U_{2\ell}$ so we have $f_j - \ell \leq Ng_j \leq f_j$ for $j \in \{1, 2\}$. Since $\ell < \gamma\sqrt{\ell}N$, it suffices to prove that for every real number y it holds that

$$\Pr[y - 2\gamma\sqrt{\ell} \leq g_j \leq y + 2\gamma\sqrt{\ell}] \leq e^{-4C} .$$

As g_j is an Irwin-Hall distribution (i.e. the sum of i.i.d. copies of $U[0, 1]$) with mean $\ell/2$, the maximum of the left hand side is obtained when $y = \ell/2$ so it suffices to prove that

$$\Pr\left[\frac{\ell}{2} - 2\gamma\sqrt{\ell} \leq g_j \leq \frac{\ell}{2} + 2\gamma\sqrt{\ell}\right] \leq e^{-4C} . \quad (4)$$

Since $U_i \sim U[0, 1]$, it has mean $1/2$ and standard deviation $1/\sqrt{12}$, so we have by the Central Limit Theorem that

$$\lim_{\ell \rightarrow \infty} \Pr\left[\frac{\ell}{2} - 2\gamma\sqrt{\ell} \leq g_j \leq \frac{\ell}{2} + 2\gamma\sqrt{\ell}\right] = \Phi(2\sqrt{12}\gamma) - \Phi(-2\sqrt{12}\gamma) = 2\Phi(2\sqrt{12}\gamma) - 1 .$$

Now, choose γ such that $2\Phi(2\sqrt{12}\gamma) - 1 = e^{-4C}/2$. Then we have

$$\lim_{\ell \rightarrow \infty} \Pr\left[\frac{\ell}{2} - 2\gamma\sqrt{\ell} \leq g_j \leq \frac{\ell}{2} + 2\gamma\sqrt{\ell}\right] = \frac{1}{2e^{4C}} ,$$

implying that for all ℓ sufficiently large as a function of C we have that (4) holds. \square

Claim 3. *Let $j \in \{1, 2\}$. It holds that*

$$\Pr[|f_j - \ell(N+1)/2| \geq \ell^{2/3}N] \leq \frac{1}{\ell} . \quad (5)$$

Proof. As in the proof of the previous claim, since $f_j - \ell \leq Ng_j \leq f_j$ and since $2\ell \leq \ell^{2/3}N$, it suffices to prove that

$$\Pr\left[|g_j - \ell/2| \geq \frac{1}{2}\ell^{2/3}\right] \leq \frac{1}{\ell} .$$

Since g_j is the sum of ℓ i.i.d. random variables, each in $[0, 1]$, and each with mean $\frac{1}{2}$, it follows by

Chernoff's inequality (see, e.g. [2], Appendix A) that

$$\Pr \left[|g_i - \ell/2| \geq \frac{1}{2} \ell^{2/3} \right] \leq 2e^{-\ell^{4/3}/(8\ell)} = 2e^{-\ell^{1/3}/8} \leq \frac{1}{\ell}.$$

□

Armed with the three claims we proceed as follows. Since f_1 and f_2 are independent and since $\ell \leq f_j \leq \ell N$ we have that, for any t , $0 \leq t \leq 2\ell N$,

$$\Pr[f_1 + f_2 = t] = \sum_{k=\ell}^{\ell N} \Pr[f_1 = k] \cdot \Pr[f_2 = t - k].$$

We cover $\{\ell, \dots, \ell N\}$ with five (not necessarily disjoint) sets S_1, S_2, S_3, S_4, S_5 defined as follows.

$$\begin{aligned} S_1 &= \{k \mid \ell(N+1)/2 - \gamma\sqrt{\ell}N \leq k \leq \ell(N+1)/2 + \gamma\sqrt{\ell}N\} \\ S_2 &= \{k \mid \ell(N+1)/2 - \gamma\sqrt{\ell}N \leq t - k \leq \ell(N+1)/2 + \gamma\sqrt{\ell}N\} \\ S_3 &= \{k \mid |k - \ell(N+1)/2| \geq \ell^{2/3}N\} \\ S_4 &= \{k \mid |(t - k) - \ell(N+1)/2| \geq \ell^{2/3}N\} \\ S_5 &= \{\ell, \dots, \ell N\} \setminus (S_1 \cup S_2 \cup S_3 \cup S_4). \end{aligned}$$

For $z \in \{1, 2, 3, 4, 5\}$ let $J_z = \sum_{k \in S_z} \Pr[f_1 = k] \cdot \Pr[f_2 = t - k]$ so that we have

$$\Pr[f_1 + f_2 = t] \leq J_1 + J_2 + J_3 + J_4 + J_5.$$

We now bound each J_z where we will use Claim 1, Claim 2, Claim 3, and the trivial bound $\Pr[f_j = k'] \leq 1/N$ which holds for every $k' \in [N]$ since f_j is the sum of discrete random variables, each uniform on N possible values. By the definition of S_1 and by Claim 2 applied to f_1 with $x = \ell(N+1)/2$:

$$J_1 = \sum_{k \in S_1} \Pr[f_1 = k] \cdot \Pr[f_2 = t - k] \leq \frac{1}{N} \sum_{k \in S_1} \Pr[f_1 = k] \leq \frac{1}{Ne^{4C}}.$$

Similarly, by the definition of S_2 and by Claim 2 applied to f_2 with $x = \ell(N+1)/2$:

$$J_2 = \sum_{k \in S_2} \Pr[f_1 = k] \cdot \Pr[f_2 = t - k] \leq \frac{1}{N} \sum_{k \in S_2} \Pr[f_2 = t - k] \leq \frac{1}{Ne^{4C}}.$$

By the definition of S_3 and by Claim 3 applied to f_1 :

$$J_3 = \sum_{k \in S_3} \Pr[f_1 = k] \cdot \Pr[f_2 = t - k] \leq \frac{1}{N} \sum_{k \in S_3} \Pr[f_1 = k] \leq \frac{1}{N\ell} \leq \frac{1}{Ne^{4C}}.$$

By the definition of S_4 and by Claim 3 applied to f_2 :

$$J_4 = \sum_{k \in S_4} \Pr[f_1 = k] \cdot \Pr[f_2 = t - k] \leq \frac{1}{N} \sum_{k \in S_4} \Pr[f_2 = t - k] \leq \frac{1}{N\ell} \leq \frac{1}{Ne^{4C}}.$$

Finally consider J_5 . To estimate it, we will distinguish between two cases, according to the value of t . Assume first that $t \leq 2\ell(N+1)/3$ or $t \geq 4\ell(N+1)/3$. In this case $S_3 \cup S_4 = \{\ell, \dots, \ell N\}$ and hence $S_5 = \emptyset$ implying that $J_5 = 0$. Assume next that $2\ell(N+1)/3 < t < 4\ell(N+1)/3$. First, observe that the number of elements of S_5 is at most $2\ell^{2/3}N + 1 < 3\ell^{2/3}N$ as it is disjoint from, say, S_3 . Consider some term of J_5 , namely $\Pr[f_1 = k] \cdot \Pr[f_2 = t - k]$ where $k \in S_5$. By Claim 1, we have that $\Pr[f_1 = k] \leq \Pr[f_1 = k^*]$ where $k^* \in S_1$ as k^* is closer to the mean $\ell(N+1)/2$ than k is. But the number of elements in S_1 is at least $2\gamma\sqrt{\ell}N$ so we must have $\Pr[f_1 = k] \leq 1/|S_1| \leq 1/(2\gamma\sqrt{\ell}N)$. Similarly, by Claim 1, we have that $\Pr[f_2 = t - k] \leq \Pr[f_2 = t - k^*]$ where $k^* \in S_2$ as $t - k^*$ is closer to the mean $\ell(N+1)/2$ than $t - k$ is. But the number of elements in S_2 is at least $2\gamma\sqrt{\ell}N$ so we must have $\Pr[f_2 = t - k] \leq 1/|S_2| \leq 1/(2\gamma\sqrt{\ell}N)$. Hence, in any case,

$$J_5 = \sum_{k \in S_5} \Pr[f_1 = k] \cdot \Pr[f_2 = t - k] \leq 3\ell^{2/3}N \cdot \left(\frac{1}{2\gamma\sqrt{\ell}N}\right)^2 \leq \frac{1}{Ne^{4C}}.$$

We have thus proved that $\Pr[f_1 + f_2 = t] \leq 5e^{-4C}/N$, as required. \square

3 Proof of the main result

3.1 Proof of the lower bound of Theorem 1

Let ϵ be given, $0 < \epsilon < 1$. Let n be sufficiently large and m be given such that $n \leq m \leq n^{1/\epsilon}$. We shall construct a hypergraph H on n vertices and m hyperedges such that $s(H) \geq m^{2-\epsilon}$. Let r be a positive integer such that $\epsilon > 2/(r+1)$. Recall that Lemma 2 implies, for sufficiently large N and any positive δ , the existence of a hypergraph H' on N vertices and $N^{(r+1)/2+\delta}$ hyperedges satisfying $s(H') \geq cN^r$, for a constant $c = c(r)$.² Note that $N^{(r+1)/2+\delta}$ is slightly larger than the expression for the number of hyperedges given in Lemma 2, but we can always add hyperedges if necessary as this does not decrease the parameter s . Next, we choose N such that $m = N^{(r+1)/2+\delta}$. Thus H' has m hyperedges. Note that:

$$N \leq (N^{(r+1)/2+\delta})^{2/(r+1)} = m^{2/(r+1)} \leq (n^{1/\epsilon})^{2/(r+1)} \leq n,$$

so by just adding $n - N$ isolated vertices to H' we obtain a hypergraph H with n vertices and m hyperedges and with $s(H) = s(H') \geq cN^r$. Hence for δ sufficiently small

$$s(H) \geq cN^r \geq c(m^{2/(r+1+2\delta)})^r \geq m^{2-\epsilon}.$$

\square

²We ignore rounding issues as these have no effect on the asymptotic statement of the theorem.

3.2 Proof of the upper bound of Theorem 1

Consider a hypergraph $H = (V, E)$ on n vertices and m hyperedges. We shall argue that an appropriate random labeling is distinguishing with positive probability. Before we prove our upper bound m^2/C on $s(H)$, we shall quickly remark that the upper bound $s(H) \leq m^2$ is easy to obtain. Indeed, assign an integer value from $[m^2]$ to each vertex independently with probability $1/m^2$. Consider the probability p that two given distinct hyperedges e and e' get the same sum of the labels. Fix an arbitrary vertex y in the symmetric difference of e and e' . Then assuming that all other labels in the union of e and e' are fixed, there is at most one value of the label assigned to y that makes the sum of labels in e and e' the same. Thus $p \leq 1/m^2$. Taking the union bound over all $\binom{m}{2}$ pairs of hyperedges, we see that the probability that the labeling is not distinguishing is at most $\binom{m}{2}/m^2 < 1$.

Next we shall improve this easy upper bound to $s(n, m) = o(m^2)$. This turns out to require significantly more effort. We first describe the main idea of the proof. We consider a hypergraph $H = (V, E)$ on n vertices and m hyperedges. Let $C > 0$ and $N = \lceil m^2/C \rceil$. Consider a labeling $f : V(H) \rightarrow [N]$ such that $f(v)$ is assigned randomly with $\Pr[f(v) = i] = 1/N$ for any $i \in [N]$ and assignment of values to distinct vertices is independent. Let, for any set Q of vertices, $s(Q)$ denote $\sum_{v \in Q} f(v)$. For two hyperedges e, e' , let $X(e, e') = e \setminus e'$. Observe that a vertex labeling f is distinguishing on H if for any two hyperedges $s(X(e, e')) \neq s(X(e', e))$. Let $B(e, e')$ be the (bad) event that $s(X(e, e')) = s(X(e', e))$.

Consider sets $D(e, e') = X(e, e') \cup X(e', e)$ and split the analysis into cases depending on the size of $D(e, e')$. For small $D(e, e')$ we would like to apply the Lovász Local Lemma, but of course the lemma's dependency digraph might have a high degree if there are vertices that belong to many such $D(e, e')$'s, called "dangerously popular" vertices. We treat them first observing that there are not so many such vertices. Finally, we deal with large $D(e, e')$'s. For those we show that the bad event $B(e, e')$ does not happen by choosing a large set S of size 2ℓ in $X(e, e')$ or in $X(e', e)$, fixing the labels on the remaining vertices in $D(e, e')$ and showing that $\Pr[B(e, e')] \leq \Pr[s(S) = t]$ for a specific value t , finally upper-bounding the latter using Lemma 3. We now proceed with the detailed proof.

For our fixed C , let $K > P > C$ where K and P are positive integer constants chosen to satisfy the claimed inequalities used in the proof. They will only depend on C . For the rest of the proof we assume that $C > 3$ and note that if the theorem holds for some value of C , it holds for any smaller positive value of C .

- A pair of hyperedges e, e' is *dangerous* if $|D(e, e')| \leq K$. Otherwise, the pair is called *non-dangerous*.
- We call a vertex $w \in V(H)$ *dangerously popular* if for at least m^2/K^3 dangerous pairs e, e' it holds that $w \in D(e, e')$. Let S be the set of all dangerously popular vertices.
- For a pair $e, e' \in E(H)$ (whether dangerous or not) let $Y(e, e') = X(e, e') \cap S$, the set of dangerously popular vertices in $X(e, e')$ and let $Z(e, e') = X(e, e') \setminus Y(e, e')$.

- We call a pair $e, e' \in E(H)$ *special* if each vertex of $D(e, e')$ is dangerously popular, i.e. $D(e, e') = Y(e, e') \cup Y(e', e)$.
- Two special pairs e_1, e'_1 and e_2, e'_2 are *equivalent* if $\{X(e_1, e'_1), X(e'_1, e_1)\} = \{X(e_2, e'_2), X(e'_2, e_2)\}$. Observe that “equivalent” is an equivalence relation over the special pairs.
- We call a non-dangerous and non-special pair $e, e' \in E(H)$ *newly dangerous* if all but at most P vertices of $D(e, e')$ are dangerously popular (so $1 \leq |Z(e, e') \cup Z(e', e)| \leq P$ for such pairs).

We observe that that the number of dangerously popular vertices is $|S| \leq K^4$. Indeed, the total sum of cardinalities of all the $D(e, e')$'s ranging over all dangerous pairs is at most $K \binom{m}{2}$ and as each dangerously popular vertex is counted at least m^2/K^3 times, there are at most $K \binom{m}{2} / (m^2/K^3) \leq K^4$ dangerously popular vertices.

Recall that $N = \lceil m^2/C \rceil$. Our assignment of values from $[N]$ to the vertices of H proceeds in two steps. We will first assign values to the dangerously popular vertices such that some properties are guaranteed. We will then assign values to the remaining vertices.

Step 1: Assign random values to the dangerously popular vertices (i.e. the vertices in S). As in the proof of Lemma 3, for the purpose of our analysis, the random values are assigned as follows. Each $w \in S$ is assigned uniformly and independently a random *real* $g(w)$ in $[0, N]$. Then, we define $f(w) = \lceil g(w) \rceil$. Since $\Pr[f(w) = 0] = 0$, we have that $f(w)$ is discrete uniform in $[N]$ as the probability that $f(w) = t$ is $1/N$ for each $t \in [N]$.

Recall that $Y(e, e') = X(e, e') \cap S$. Let $f(e, e') = \sum_{w \in Y(e, e')} f(w)$. We say that Step 1 is *successful* if both of the following hold:

1. For every special pair e, e' we have $f(e, e') \neq f(e', e)$.
2. For at most $m^2 e^{-4C}$ newly dangerous pairs e, e' it holds that $|f(e, e') - f(e', e)| \leq PN$.

Lemma 4. *With positive probability, Step 1 is successful.*

Lemma 4 will be proved later, but for now assume that it holds, so fix an assignment of the vertices of S such that Step 1 is successful.

Step 2: Assign random values to the remaining $n - |S|$ vertices. As in Step 1, we assign the random values as follows. Each $w \in V(H) \setminus S$ is assigned uniformly and independently a random *real* $g(w)$ in $[0, N]$. Then, we define $f(w) = \lceil g(w) \rceil$. Recall that $f(w)$ is discrete uniform in $[N]$. This now defines for each hyperedge $e \in E(H)$ the sum $s(e) = \sum_{w \in e} f(w)$. We need to estimate the probability that $s(e) = s(e')$ for distinct hyperedges e, e' . We partition the pairs (e, e') of hyperedges into five types:

- (a) The special pairs.
- (b) The newly dangerous pairs for which $|f(e, e') - f(e', e)| > PN$.
- (c) The newly dangerous pairs for which $|f(e, e') - f(e', e)| \leq PN$.
- (d) Non-dangerous pairs that are not newly dangerous and not special.

(e) Dangerous pairs that are not special.

We refer to these types by their letter. Each pair of hyperedges is of precisely one of these types. We now analyze each type. Let A_a, A_b, A_c, A_d , and A_e be events that there is a pair e, e' of type (a), (b), (c), (d), or (e), respectively, such that $s(e) = s(e')$. We prove the following lemmas later.

Lemma 5. $\Pr[A_a] = \Pr[A_b] = 0$.

Lemma 6. $\Pr[A_c] \leq e^{-3C}$.

Lemma 7. $\Pr[A_d] \leq e^{-3C}$.

Lemma 8. $\Pr[A_e] \leq 1 - e^{-2C}$.

Lemmas 5, 6, 7, and 8 imply that $\Pr[A_a \cup A_b \cup A_c \cup A_d \cup A_e] \leq e^{-3C} + e^{-3C} + 1 - e^{-2C} < 1$. Thus, with positive probability none of these bad events happen and there is a desired distinguishing labeling of H . It remains to prove Lemmas 4, 5, 6, 7, and 8.

In several proofs we shall need the following observation for any distinct subsets X and X' of vertices, recalling that $s(X) = \sum_{w \in X} f(w)$,

$$\Pr(s(X) = s(X')) \leq 1/N. \quad (6)$$

The reason for this observation to hold is the same as we outlined in the first paragraph of the proof - fixing all labels except for one vertex, say y , in the symmetric difference of X and X' , we see that $\Pr(s(X) = s(X')) \leq \Pr(f(y) = t) = 1/N$, for some specific value t .

Proof of Lemma 5. If e, e' is a pair of type (a), then clearly $s(e) - s(e') = f(e, e') - f(e', e)$. But since Step 1 is successful, we have that $f(e, e') \neq f(e', e)$ and hence $s(e) \neq s(e')$. Thus the event A_a never happens.

If e, e' is a pair of type (b), i.e., a newly-dangerous pair for which $|f(e, e') - f(e', e)| > PN$ we proceed as follows. Assume without loss of generality that $f(e, e') - f(e', e) > PN$. Clearly

$$s(e) = f(e, e') + \sum_{w \in e \cap e'} f(w) + \sum_{w \in Z(e, e')} f(w) \geq f(e, e') + \sum_{w \in e \cap e'} f(w).$$

On the other hand,

$$s(e') = f(e', e) + \sum_{w \in e \cap e'} f(w) + \sum_{w \in Z(e', e)} f(w) \leq f(e', e) + \sum_{w \in e \cap e'} f(w) + PN,$$

because $|Z(e', e)| \leq P$ by the definition of newly-dangerous. It follows from the last two inequalities that

$$s(e) - s(e') \geq f(e, e') - f(e', e) - PN > 0,$$

so we have that $s(e) \neq s(e')$. Thus the event A_b never happens. \square

Proof of Lemma 6. Let e, e' be a pair of type (c), namely a newly dangerous pair for which it holds that $|f(e, e') - f(e', e)| \leq PN$. As Step 1 is successful, we have that the number of pairs of type (c) is at most $m^2 e^{-4C}$.

By (6), we have that $\Pr[s(e) = s(e')] \leq 1/N$. Since the number of pairs of type (c) is at most $m^2 e^{-4C}$ we have that

$$\Pr[A_c] \leq \frac{m^2 e^{-4C}}{N} = \frac{m^2 e^{-4C}}{\lfloor m^2/C \rfloor} \leq C e^{-4C} \leq e^{-3C}.$$

□

Proof of Lemma 8. For a pair e, e' of type (e), let $A(e, e')$ be the event that $s(e) = s(e')$. Using (6) we have $\Pr[A(e, e')] \leq 1/N$. Letting L denote the set of pairs of type (e), our goal is to prove that $\Pr[\cap_{\{e, e'\} \in L} \overline{A(e, e')}] \geq e^{-2C}$ as this is equivalent to proving that $\Pr[A_e] \leq 1 - e^{-2C}$. To this end, we will use the Lovász Local Lemma (LLL). Consider the dependency digraph on the events $A(e, e')$ (note: there could be as many as $\binom{m}{2}$ such events). We claim that any event $A(e, e')$ depends on not too many other events. Indeed, if $Z(e_1, e'_1) \cup Z(e'_1, e_1)$ is disjoint from $Z(e_2, e'_2) \cup Z(e'_2, e_2)$, then the event $A(e_1, e'_1)$ is independent of the event $A(e_2, e'_2)$ as they involve assignment of values to disjoint sets of vertices. Recall that the pairs of type (e) are, in particular, dangerous pairs. Hence $|Z(e, e') \cup Z(e', e)| \leq K$, for any pair e, e' of type (e). Furthermore, each vertex of $Z(e, e') \cup Z(e', e)$ is not dangerously popular. Thus, we have that $A(e, e')$ is independent of all but at most $K \cdot m^2 / K^3 = m^2 / K^2$ other events. Denote $\{e_1, e'_1\} \sim \{e_2, e'_2\}$ if $Z(e_1, e'_1) \cup Z(e'_1, e_1)$ is not disjoint from $Z(e_2, e'_2) \cup Z(e'_2, e_2)$. To apply LLL, define $x(e, e') = 2/N$. For any e_1, e_2 of type (e) it now holds that

$$\begin{aligned} x(e_1, e_2) \prod_{\{e'_1, e'_2\} \sim \{e_1, e_2\}} (1 - x(e'_1, e'_2)) &\geq \frac{2}{N} \left(1 - \frac{2}{N}\right)^{m^2/K^2} \\ &\geq \frac{2}{N} \left(1 - \frac{2}{N}\right)^{CN/K^2} > \frac{1}{N} \geq \Pr[A(e_1, e_2)] \end{aligned}$$

so the condition in the statement of LLL holds. So, by the LLL, we have that

$$\Pr[\cap_{\{e, e'\} \in L} \overline{A(e, e')}] \geq (1 - x(e, e'))^{|L|} \geq \left(1 - \frac{2}{N}\right)^{m^2/2} \geq e^{-2C},$$

as required. □

Proof of Lemma 4. We first prove that with probability at least $2/3$, for every special pair e, e' we have $f(e, e') \neq f(e', e)$. Observe that the number of equivalence classes in the “equivalent” relation is at most $2^{|S|} 2^{|S|} \leq 4^{K^4}$ (namely, a constant). Since for two equivalent special pairs e_1, e'_1 and e_2, e'_2 we have that $f(e_1, e'_1) \neq f(e'_1, e_1)$ if and only if $f(e_2, e'_2) \neq f(e'_2, e_2)$, it suffices to consider a representative special pair from every equivalence class. Now, if e, e' is a special pair then, using (6) we have that $\Pr[f(e, e') = f(e', e)] \leq 1/N$. We have by the union bound that the probability

that for some special pair $f(e, e') = f(e', e)$ is at most $4K^4/N \ll 1/3$. So, with probability at least $2/3$, for every special pair e, e' we have $f(e, e') \neq f(e', e)$.

We next prove that with probability at least $2/3$, for at most m^2e^{-4C} newly dangerous pairs it holds that $|f(e, e') - f(e', e)| \leq PN$ (thus, we will have that Step 1 is successful with probability at least $1 - (1 - 2/3) - (1 - 2/3) > 0$, as required). To prove this we will need to establish some “anti-concentration” result, and this will be possible by applying the law of large numbers to some appropriate random variable.

Let us fix a newly dangerous pair u, v . We know that u, v is not a dangerous pair, namely $|D(e, e')| \geq K$. On the other hand, we know that $D(e, e')$ contains many dangerously popular vertices, since $1 \leq |Z(e, e')| \leq P$. So, either $|Y(e, e')| \geq (K - P)/2 \geq K/4$ or else $|Y(e', e)| \geq (K - P)/2 \geq K/4$. Assume without loss of generality that $|Y(e, e')| \geq K/4$. Now, suppose we are given that $f(e', e) = t$ for some integer t . Given this information, we would like to upper bound the probability that $f(e, e')$ lies in $[t - PN, t + PN]$. If we can provide an upper bound which does not depend on t , then we have upper-bounded the probability that $|f(e, e') - f(e', e)| \leq PN$ regardless of any given information.

So, consider indeed the random variable $f(e, e')$. It is the sum of $\ell = |Y(e, e')| \geq K/4$ i.i.d. random variables, namely $f(e, e') = X_1 + \dots + X_\ell$ where each X_i is discrete uniform in $[N]$. It will be slightly more convenient to normalize as follows. Recall that each X_i corresponds to some $f(w)$ for $w \in Y(e, e')$ and that $f(w) = \lceil g(w) \rceil$. Hence X_i is determined by first selecting uniformly at random a real number W_i in $[0, N]$ and then setting $X_i = \lceil W_i \rceil$. Define $U_i = W_i/N$ and notice that $U_i \sim U[0, 1]$ and that $X_i = \lceil NU_i \rceil$.

Let $g(e, e') = U_1 + \dots + U_\ell$ and observe that $f(e, e') - \ell \leq Ng(e, e') \leq f(e, e')$. Thus, it suffices to upper bound the probability that $g(e, e')$ lies in $[t/N - P - \ell/N, t/N + P]$. As $\ell/N \leq K/N \leq P$, it suffices to upper bound the probability that $g(e, e')$ lies in $[t^* - 2P, t^* + 2P]$ for some real number t^* . As $U_1 + \dots + U_\ell$ is an Irwin–Hall distribution which is concave in $[0, \ell]$, the latter probability is maximized when $t^* = \ell/2$, so it remains to upper bound the probability that $g(e, e')$ lies in $[\ell/2 - 2P, \ell/2 + 2P]$. As the U_i are i.i.d. each having mean $\frac{1}{2}$ and standard deviation $1/\sqrt{12}$ (i.e. absolutely bounded standard deviation), the (weak) law of large numbers applies to their sum $g(e, e')$, namely for every constant P

$$\lim_{\ell \rightarrow \infty} \Pr[g(e, e') \in [\ell/2 - 2P, \ell/2 + 2P]] = 0.$$

This, in turn, means that for all K sufficiently large as a function of P, C (hence all ℓ sufficiently large since $\ell \geq \lfloor K \rfloor / 4$),

$$\Pr[g(e, e') \in [\ell/2 - 2P, \ell/2 + 2P]] \leq \frac{1}{3}e^{-4C}.$$

We have thus proved that $\Pr[|f(e, e') - f(e', e)| \leq PN] \leq e^{-4C}/3$. As there are less than m^2 pairs to consider, we have that the expected number of newly dangerous pairs satisfying $|f(e, e') - f(e', e)| \leq PN$ is at most $m^2e^{-4C}/3$. By Markov’s inequality the probability that there are more than m^2e^{-4C} such pairs is less than $1/3$, so indeed with probability at least $2/3$, for at most m^2e^{-4C} newly dangerous pairs it holds that $|f(e, e') - f(e', e)| \leq PN$. \square

Proof of Lemma 7. Let e, e' be a pair of type (d), namely it is a non-dangerous pair and is not newly dangerous nor special. We will prove that $\Pr[s(e) = s(e')] \leq e^{-3C}/m^2$. The lemma then follows as there are less than m^2 such pairs to consider. Being non-dangerous and not newly dangerous means that $|Z(e, e') \cup Z(e', e)| \geq P$. Assume without loss of generality that $|Z(e, e')| \geq P/2$. Let $\ell = \lfloor P/4 \rfloor$ and let Z be a subset of $Z(e, e')$ of size 2ℓ .

Suppose we are given the value of $f(w)$ for all $w \in V(H) \setminus Z$. Then, conditioned on this information, for $s(e) = s(e')$ to hold, $s(Z)$ must avoid a particular value t . By Lemma 3 we have that $\Pr[s(Z) = t] \leq 5/(e^{4C}N)$. Thus using the union bound over all pairs of hyperedges of type (d), we have that

$$\Pr[A_d] \leq \frac{m^2}{2} \frac{5}{e^{4C}N} \leq \frac{5C}{e^{4C}} \leq e^{-3C} .$$

□

4 Proofs of Theorems 2, 3, 4.

4.1 Proof of Theorem 2

Let G be a nonempty n -vertex graph with maximum degree Δ , minimum degree δ , and the largest number of vertices with pairwise distinct closed neighbourhoods equal to n' . We aim to show that

$$\frac{n' + \delta}{\Delta + 1} \leq s^*(G) \leq \max\{(n - d(v) - 1)(d(v) + 1) + 2 : v \in V(G)\} .$$

For a vertex labeling f , we say that a pair of vertices u, v is *bad* if $N[u] \neq N[v]$ and $s_f^*(u) = s_f^*(v)$, otherwise the pair is *good*. Thus, a labeling is vertex sum-distinguishing if all pairs are good. Let $\xi = \max\{(n - d(v) - 1)(d(v) + 1) + 2 : v \in V(G)\}$, where $d(v)$ is the degree of vertex v . Consider a labeling $f : V(G) \rightarrow [\xi]$ with a smallest number of bad pairs. We argue that the number of bad pairs is, in fact, zero.

If not, let u, v be a bad pair. Let $x = u$ if u and v are not adjacent and otherwise let $x = w$, for some $w \in (N(v) \setminus N(u)) \cup (N(u) \setminus N(v))$. Note that changing the label for x makes the pair u, v good. We shall change the label of x such that no good pair becomes bad, i.e., so that the number of bad pairs decreases. Denote the new labeling f' . Let t be a new value assigned to x , i.e., $t \neq f(x)$, $f'(x) = t$, $f'(z) = f(z)$, for any $z \in V(G) - x$.

We see that $s_{f'}^*(y) = s_f^*(y)$ if $y \notin N[x]$ and $s_{f'}^*(y) = s_f^*(y) - f(x) + t$ if $y \in N[x]$. Thus $s_{f'}^*(y) \neq s_{f'}^*(y')$ if $s_f^*(y) \neq s_f^*(y')$ and $(y, y' \in N[x])$ or $y, y' \in V(G) - N[x]$. We have that $s_{f'}^*(y) = s_{f'}^*(y')$ for $y \in N[x]$ and $y' \notin N[x]$ if and only if $s_f^*(y) - f(x) + t = s_f^*(y')$. So, a new bad pair can only appear if one vertex is in $N[x]$ and another is not.

Choose $t \in Q$, where

$$Q = [\xi] \setminus (\{f(x)\} \cup \{s_f^*(y') - s_f^*(y) + f(x) : y \in N[x], y' \notin N[x]\}) .$$

Since

$$\begin{aligned}
|\{s_f^*(y') - s_f^*(y) + f(x) : y \in N[x], y' \notin N[x]\}| &\leq |\{(y, y') : y \in N[x], y' \notin N[x]\}| \\
&= (n - d(x) - 1)(d(x) + 1) \\
&\leq \xi - 2,
\end{aligned}$$

the set Q is non-empty, so there is a choice of $t \leq \xi$, such that $t \neq f(x)$ and $s_{f'}^*(y) \neq s_{f'}^*(y')$ for any $y \in N[x]$ and $y' \notin N[x]$. Since there is no bad pair y, y' for $y, y' \in N[x]$ or $y, y' \notin N[x]$ in f' that was not bad in f and the pair u, v that was bad in f is no longer bad in f' , we see that the number of bad pairs in f' is strictly less than the number of bad pairs in f , a contradiction.

For the lower bound, observe that if f is a vertex sum-distinguishing labeling of G with the largest label k , then $S = \{s_f^*(v) : v \in V(G)\} \subseteq \{(\delta(G) + 1) \cdot 1, \dots, (\Delta + 1) \cdot k\}$. Since $|S| \geq n'$, we have $n' \leq |S| \leq (\Delta + 1)k - (\delta + 1) + 1$, giving the desired lower bound.

Note that the lower bound is tight for any pair δ, Δ , $\delta \leq \Delta$. If $\delta = \Delta$ consider $G = K_{\Delta+1}$, for which $n' = 1$, $s^*(G) = 1$, and $(n' + \delta)/(\Delta + 1) = 1$. If $\delta < \Delta$, consider G that is a vertex-disjoint union of $K_{\Delta+1}$ and $K_{\delta+1}$. In this case $n' = 2$, $s^*(G) = 1$, and $\lceil (n' + \delta)/(\Delta + 1) \rceil = 1$. \square

4.2 Proof of Theorem 3

For a tree T and its vertex u , let $L(u)$ be the set of leaves adjacent to u . Let $L(T) = \max\{|L(u)| : u \in V(T)\}$. We shall prove for $n \geq 3$ and any tree T on n vertices, that $s^*(T) \leq 2n - 2 - L(T)$ by induction on n . Note that this bound is sharp for stars.

The case $n = 3$ holds vacuously since T is a star. Suppose the statement holds for $n \geq 3$, and suppose T has $n + 1$ vertices and is not a star. Choose a vertex u for which $L(T) = |L(u)|$, choose a leaf v adjacent to u , and let $T^* = T - v$. Then $L(T) - 1 = |L(u)| - 1 \leq L(T^*)$. By induction there is a vertex sum-distinguishing labeling $f : V(T^*) \rightarrow [2n - 2 - L(T^*)]$ of T^* . Observe that

$$2n - 2 - L(T^*) \leq 2n - 2 - (L(T) - 1) = 2n - 1 - L(T) < 2n - L(T) = 2(n + 1) - 2 - L(T).$$

We define a labeling $f' : V(T) \rightarrow [2n - L(T)]$ such that $f'(u) = f(u)$, $u \in V(T^*)$, $f'(v) = \xi$.

We argue that we can find an appropriate ξ so that the labeling f' does not contain bad pairs, i.e., pairs of vertices y, y' such that $N[y] \neq N[y']$ but $s_{f'}^*(y) = s_{f'}^*(y')$. Since there are no bad pairs in T^* under f , we have that y, y' is not a bad pair if $y, y' \in T - \{u, v\}$. Thus we need to consider only the pairs y, y' , where $y \in \{u, v\}$. Let $L = L(u)$ in T .

If $y = u$ and $y' \in L$, we see that $s_{f'}^*(u) > s_{f'}^*(y')$ regardless of ξ . Thus such a pair y, y' is not bad. A pair $y = u, y' = u'$, $u' \in V(T^*) - (\{u\} \cup L)$ can be bad if $s_{f'}^*(u) = s_f^*(u) + \xi = s_f^*(u')$. A pair $y = v, y' = u'$, $u' \in V(T^*) - u$ can be bad if $s_{f'}^*(v) = f(u) + \xi = s_f^*(u')$. Thus, if $\xi \notin X$, where

$$X = \{s_f^*(u') - s_f^*(u) : u' \in V(T^*) - (\{u\} \cup L)\} \cup \{s_f^*(u') - f(u) : u' \in V(T^*) - u\},$$

then f' has no bad pairs on T . Note that $|X| \leq (n - 1) - |L| + (n - 1) = 2n - |L| - 2$. Thus, there is an available choice for ξ in $[2n - |L|] - X$. \square

4.3 Proof of Theorem 4

Let H be the hypergraph from Lemma 2 and let H^* be the dual hypergraph of H , so $\text{irr}(H^*) = s(H)$. Let n be the number of vertices of H^* that is the number of edges in H , i.e. $n = |E(H)| = \Theta(N^{(r+1)/2} \sqrt{\log N})$. We also have that $s(H) = cN^r$, so for $\epsilon < 2/(r+1)$ we have

$$\text{irr}(H^*) \geq cN^r \geq Cn^{2r/(r+1)}/\log^r n \geq n^{2-\epsilon}.$$

□

5 Concluding remarks and open problems

As mentioned in the introduction, there are connections between the considered problem and Sidon sets. Recall that a $B_h[1]$ -Sidon set is a set X of integers such that for any integer q , there is at most one subset X' of X , $|X'| = h$, so that the sum of elements from X' is q . The following bounds on the sizes of Sidon sets are known: if $X \subseteq [K]$ and X is a $B_h[1]$ -Sidon, then $|X| \leq (h \cdot h!K)^{1/h}(1 + o(1))$, see for example [14, 15]. Let $c(h)$ be a constant depending on h only such that $|X| \leq (c(h)K)^{1/h}(1 + o(1))$ for any $B_h[1]$ -Sidon set X , $X \subseteq [K]$. Consider a hypergraph H that is a union of the complete h -uniform hypergraph on N vertices and $\binom{N}{h} - N$ isolated vertices. Then H has the same number $n = N^h/h!(1 + o(1))$ of vertices and edges and $s(H) \geq (1/c(h))^h N^h = \Theta(n)$. So, this only gives a linear lower bound on $s(n) = s(n, n)$, much weaker than Theorem 1.

Note that a similar problem defined on open neighbourhoods of the vertices of a graph is equivalent to the setting we considered on the complement \overline{G} of the graph G . Indeed, if f is a vertex sum-distinguishing labeling of G , then the numbers $\sum_{u \notin N[v]} f(u)$, $v \in V(G)$ are distinct for any two vertices with distinct open neighbourhoods. We see that $V(G) - N[v] = N_{\overline{G}}(v)$, thus the sums considered correspond to the sums over open neighbourhoods in the complement.

Yet another variant of $s(H)$ is its restriction to *injective* labelings. Denoting the corresponding parameters by $s_{inj}(H)$ and $s_{inj}(n, m)$ we see, by definition, that $s_{inj}(H) \geq s(H)$ so $s_{inj}(n, m) \geq s(n, m)$. If a hypergraph H' is a union of H and all hyperedges consisting of exactly one vertex of H , then $s_{inj}(H) \leq s(H')$, and H' has at most $|E(H)| + |V(H)|$ edges. Thus, Theorem 1 trivially extends to $s_{inj}(n, m)$. By defining $s_{inj}^*(n)$ similarly and following the steps of Theorem 2, one can also show that $s_{inj}^*(G) \leq (\Delta + 2)n$, for any graph G on n vertices and maximum degree Δ .

In this paper, we addressed hypergraphs on n vertices and m hyperedges, for $n \leq m \leq n^{O(1)}$. It may be of some interest to determine the behavior of $s(n, m)$ when m is larger than a polynomial function of n . As mentioned in the introduction, the closer m gets to 2^n , the closer $s(n, m)$ gets to be linear in m .

Finally, it may be of some interest to improve the upper bound $s^*(G) \leq n(\Delta + 1)$ given in Theorem 2 for regimes of Δ that are significantly less than quadratic.

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